

1 Tropical posynomial systems

A tropical posynomial is a supremum of finitely many affine functions

$$P_{trop}(x) = \max_{a \in A} (\langle a, x \rangle + r_a)$$

where $x \in \mathbb{R}^n$ and $A \subset \mathbb{R}^n$ is a finite set (the tropical exponents).

Terminology comes from the tropical semifield $\mathbb{R}_{\max} = (\mathbb{R} \cup \{-\infty\}, \max, +)$ in which P_{trop} can be thought of as the analog of a posynomial.

For tropical posynomials Q_1, \dots, Q_n , let $S_{\mathbb{R}_{\max}}$ be the **tropical posynomial system**:

$$\forall i \in [n] \quad Q_i(x) = 0. \quad (S_{\mathbb{R}_{\max}})$$

The corresponding feasibility problem is NP-complete. These systems are the analog of real posynomial systems:

$$\forall i \in [n] \quad \sum_{a \in A_i} r_a x^a = 1, \quad (S_{\mathbb{R}})$$

where $(A_i)_{1 \leq i \leq n} \subset \mathbb{R}^n$ are finite and the r_a are positive real numbers.

3 The colorful interior of a family of convex bodies

Let $\mathcal{V} = (V_1, V_2, \dots, V_n)$ be a family of subsets of \mathbb{R}^n , each one being assigned a color.

Definition. We say a vector y is **rainbow** with respect to \mathcal{V} if $y \in \text{cone}(\cup_i V_i)$ and all linear decompositions of y use all the colors:

$$\forall \mu \in (\mathbb{R}^+)^{\cup_i V_i} \quad y = \sum_{i=1}^n \sum_{v \in V_i} \mu_v v \implies \forall i \in [n] \quad \exists v \in V_i \quad \mu_v > 0.$$

The set of rainbow vectors is denoted by $\mathfrak{m}\mathcal{V}$ and called the **colorful interior** of \mathcal{V} .

Proposition. For all $i \in [n]$, let us denote by $\mathcal{P}_i = \text{cone}(V_i)$, $\widehat{\mathcal{P}}_i = \text{cone}(\cup_{j \neq i} V_j)$ and $\mathcal{P} = \text{cone}(\cup_j V_j)$. Then

$$\mathfrak{m}\mathcal{V} = \mathcal{P} \setminus \bigcup_{i \in [n]} \widehat{\mathcal{P}}_i$$

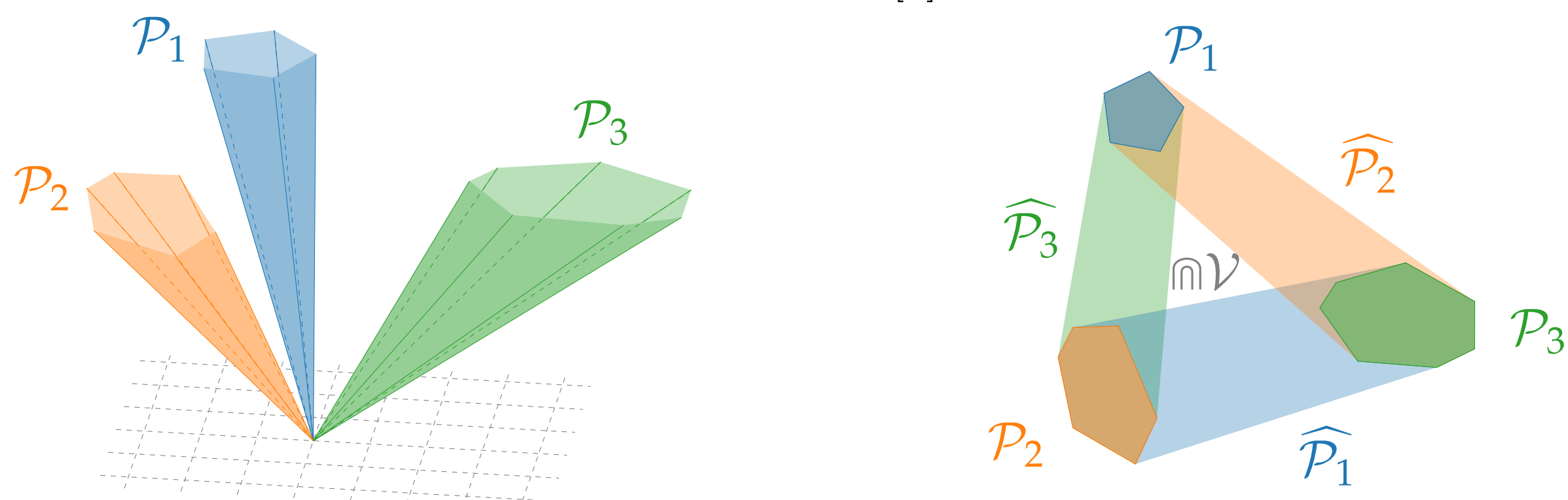


Figure 2: Three subsets V_1, V_2 and V_3 of \mathbb{R}^3 and their conic hulls (left) together with a cross-section (right). Here $\mathfrak{m}\mathcal{V}$ is the white triangle.

Definition. We say the configuration \mathcal{V} is **pointed** if $\cup_i V_i$ is contained in an open halfspace.

5 Properties of the colorful interior

We assume in this part that \mathcal{V} is a pointed configuration.

Theorem. $\mathfrak{m}\mathcal{V}$ is either empty or the interior of a simplicial cone.

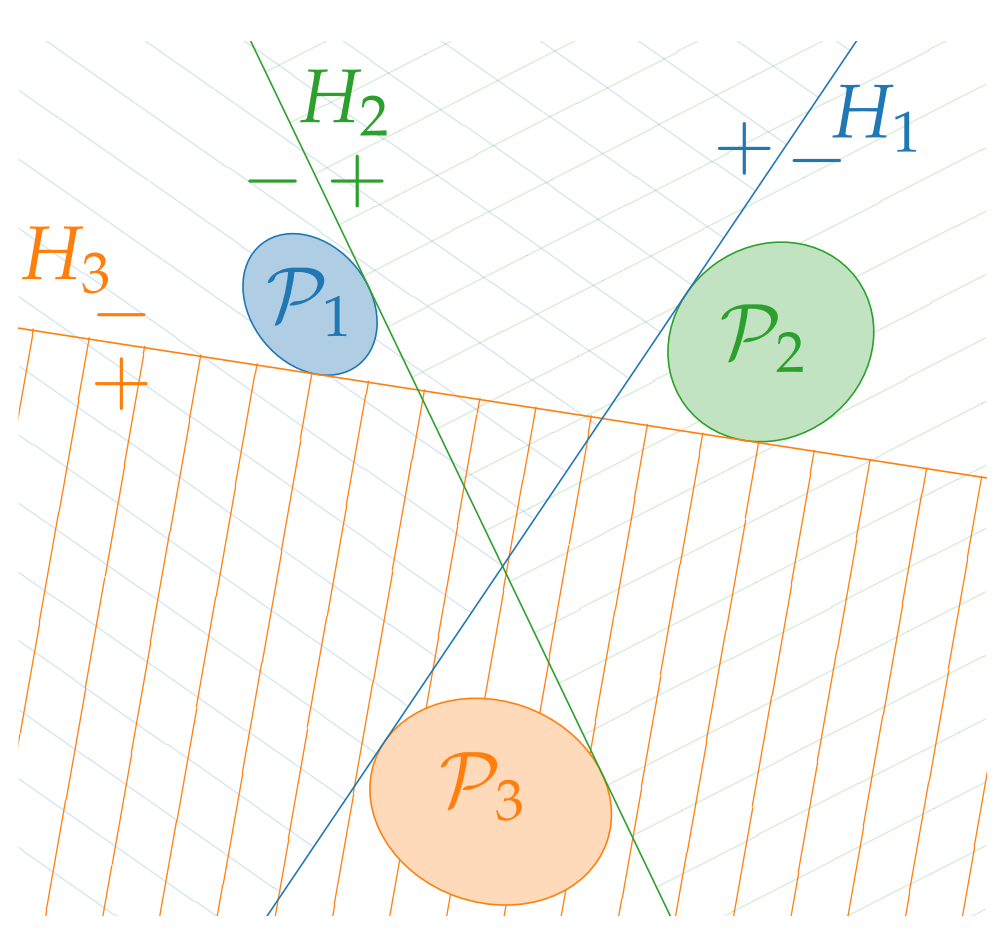


Figure 3: For convex V_1, V_2 and $V_3 \subset \mathbb{R}^3$, we have $\mathfrak{m}\mathcal{V} = H_1^{>0} \cap H_2^{>0} \cap H_3^{>0}$.

This simplicial cone is the intersection of n open halfspaces, each one

- being tangent to $n - 1$ bodies of \mathcal{V} ,
- containing the last one.

In [6], Cappell *et al.* have stated when such hyperplanes exist:

Theorem [6]. Let $\mathcal{A} = (A_1, \dots, A_n)$ be a family of n separated compact convex bodies of \mathbb{R}^n , then there are exactly two oriented affine hyperplanes that are tangent and outer to \mathcal{A} .

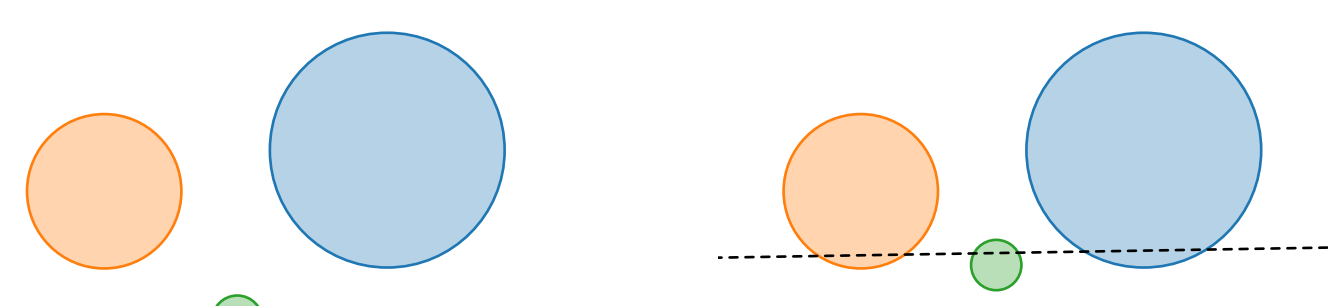


Figure 4: The three bodies are separated on the left figure, but not on the right one

For all $i \in [n]$, we define $\overline{\mathcal{P}}_i = \cap_{j \neq i} \widehat{\mathcal{P}}_j$.

Conjecture. The following assertions are equivalent:

- $\mathfrak{m}\mathcal{V} \neq \emptyset$
- The family $(\overline{\mathcal{P}}_i)_i$ is separated

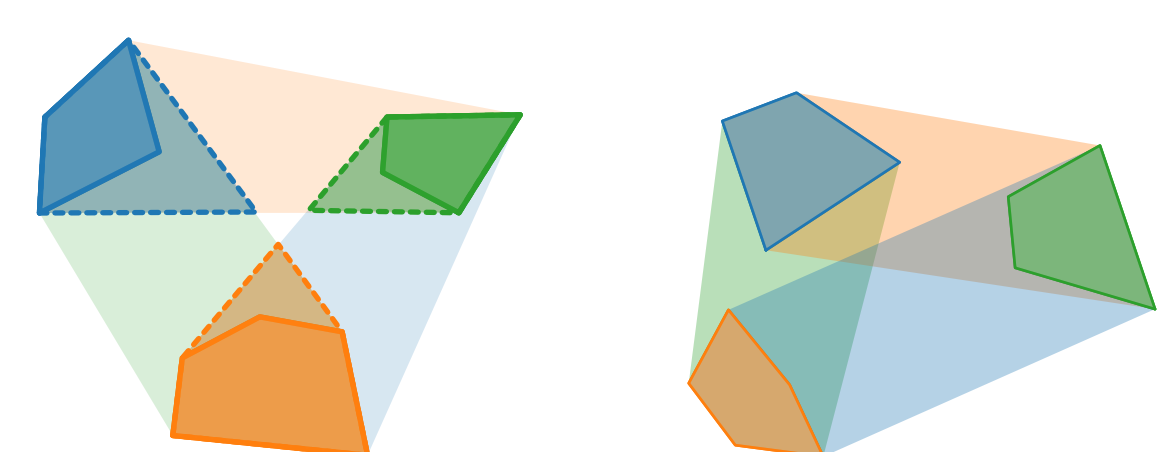


Figure 5: Left: The $(\overline{\mathcal{P}}_i)_i$ bodies visualised. Right: $\mathfrak{m}\mathcal{V} = \emptyset$ despite $(\mathcal{P}_i)_i$ are separated

We proved (i) \implies (ii) and since for all $i \in [n]$, $\mathcal{P}_i \subset \overline{\mathcal{P}}_i$, the separation of $(\mathcal{P}_i)_i$ is necessary to have $\mathfrak{m}\mathcal{V} \neq \emptyset$ (but not sufficient). We proved (ii) \implies (i) for $n = 3$.

7 References

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2 From performance evaluation to tropical posynomials

Since 2014, we collaborate with the Parisian Fire Brigade and *Préfecture de Police* to analyze the performance of their joint **emergency call center**, implementing a bi-level filtering mechanism to prioritize urgent calls.



Figure 1: Paris' 17-18-112 emergency call center

We have modeled the call center using a generalization of Petri nets [1] in order to take into account **synchronization** and **concurrency** phenomena.

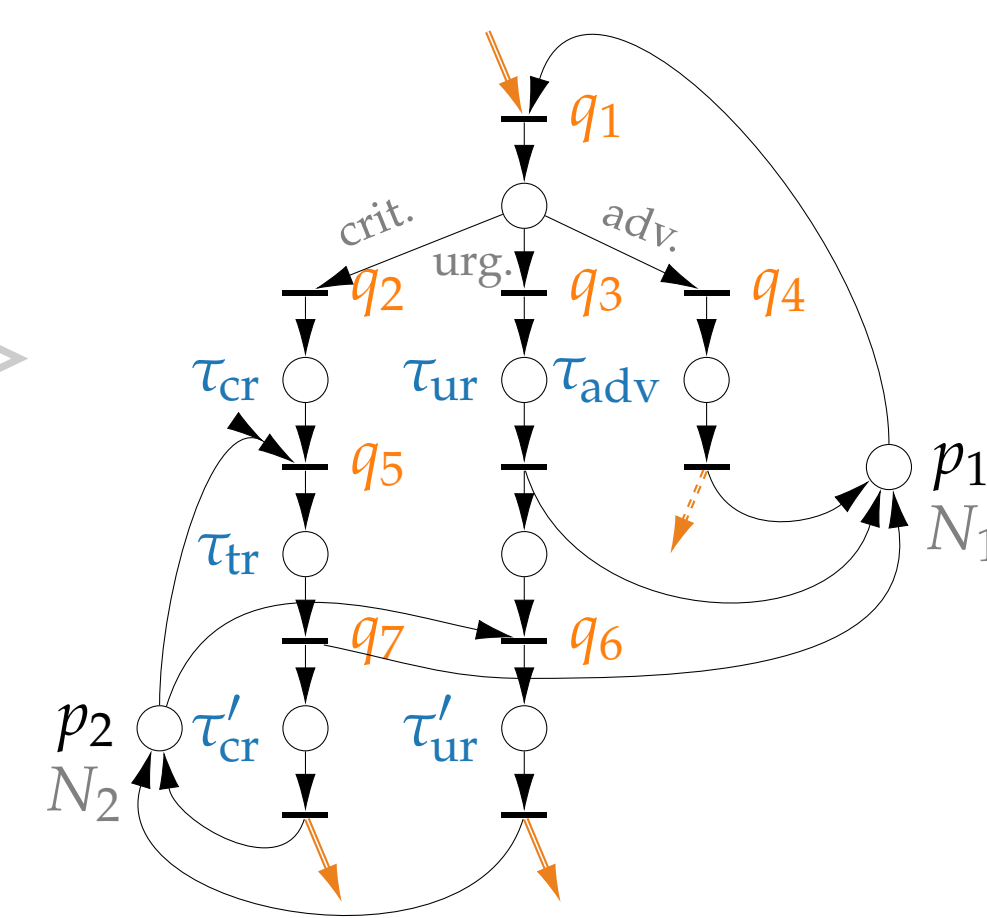


Figure 2: A Petri net model of the call center from [1]

Petri nets provide dynamic equations satisfied by **counter variables** (e.g. counting the number of calls):

$$\begin{aligned} z_1(t) &= N_1 + z_5(t - \tau_{tr}) + \pi_u z_1(t - \tau_{ur}) + \pi_{adv} z_1(t - \tau_{adv}) \\ z_5(t) &= (N_2 + z_5(t - \tau_{tr} - \tau'_{cr}) + z_6(t - \tau'_{ur}) - z_6(t^-)) \wedge \pi_{cr} z_1(t - \tau_{cr}) \\ z_6(t) &= (N_2 + z_5(t - \tau_{tr} - \tau'_{cr}) + z_6(t - \tau'_{ur}) - z_5(t)) \wedge \pi_u z_1(t - \tau_{ur}) \end{aligned}$$

Computing **stationary regimes** (i.e. $z_i(t) = \rho_i t + u_i$) reduces to a tropical posynomial system over a semifield of germs of affine functions. This builds on a series of work representing several classes of Petri nets by tropical dynamical systems [2, 3, 1].

4 Using rainbow vectors to solve tropical and real posynomial systems

Recall the tropical posynomial system we want to solve:

$$\forall i \in [n] \quad \max_{a \in A_i} (\langle a, x \rangle + r_a) = 0. \quad (S_{\mathbb{R}_{\max}})$$

We fix a vector $y \in \mathbb{R}^n$ and introduce the following linear program:

$$\min_x \langle y, x \rangle \quad \text{s.t.} \quad \forall i \in [n] \quad \max_{a \in A_i} \langle a, x \rangle + r_a \leq 0, \quad (LP)$$

with dual:

$$\sup_{\mu \geq 0} \langle \mu, r \rangle \quad \text{s.t.} \quad -y = \sum_{i=1}^n \sum_{a \in A_i} \mu_a a. \quad (LD)$$

Theorem. If (LP) is feasible and $-y$ is rainbow with respect to the configuration (A_1, \dots, A_n) , then the optimal solutions of (LP) are precisely the solutions of the tropical posynomial system.

Remark that (A_1, \dots, A_n) being a pointed configuration ensures (LP) is feasible. This theorem carries over to the real case:

Theorem. Let $S_{\mathbb{R}}$ be the following posynomial system:

$$\forall i \in [n] \quad \sum_{a \in A_i} r_a x^a = 1, \quad (S_{\mathbb{R}})$$

where $(A_i)_i$ are finite subsets of \mathbb{R}^n and $(r_a)_a$ are positive.

If (A_1, \dots, A_n) is a pointed configuration and has a non-empty colorful interior, then $S_{\mathbb{R}}$ has at least one positive solution (i.e. in $\mathbb{R}_{>0}^n$).

Proof uses the entropic program $\min_X \langle y, X \rangle \quad \text{s.t.} \quad \forall i \in [n] \quad \log \left(\sum_{a \in A_i} r_a e^{\langle a, X \rangle} \right) \leq 0$.

6 From the tropical colorful interior to tropical SVM

The definitions of the colorful interior of n convex bodies of \mathbb{R}^n can be transposed to n **tropical convex** bodies of $(\mathbb{R}_{\max})^n$.

In [4], Gärtner and Jaggi introduced the notion of **tropical support vector machines**, i.e. tropical hyperplane H separating n data point clouds, each class lying in one unique sector of H .

Theorem. The tropical colorful interior of n tropical convex bodies of $(\mathbb{R}_{\max})^n$ is the locus of all apices of tropical SVM separating these bodies. If non empty, it is a **polytrope** (i.e. it is convex tropically and classically).

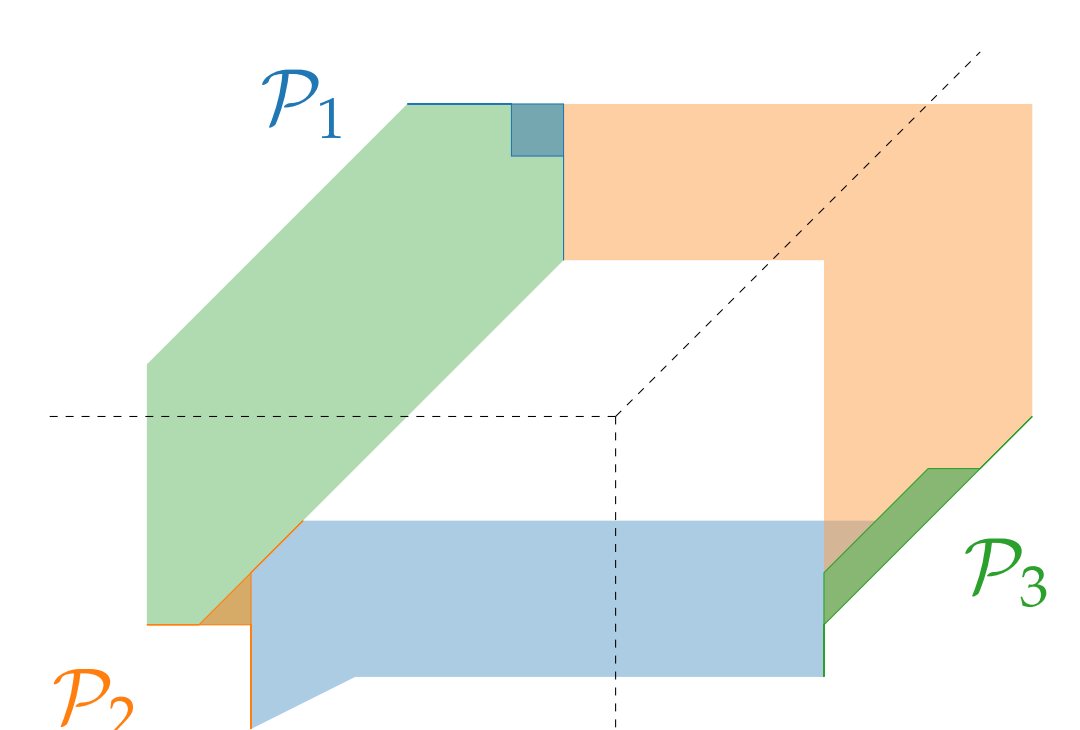


Figure 6: Three tropical simplices of $(\mathbb{R}_{\max})^3$, their colorful interior (in white), and a tropical SVM

Theorem. Deciding whether the tropical colorful interior $\mathfrak{m}\mathcal{V}$ is empty or not (and finding a point inside) can be achieved in strongly polynomial time.

This answers one of the questions raised in [4]. A key element of the proof is the following theorem inspired by a result from Lawrence and Soltan (see [5]):

Theorem. The tropical colorful interior $\mathfrak{m}\mathcal{V}$ coincides with the intersection of the interior of all multicolor tropical simplices, i.e. having one vertex in each color set:

$$\mathfrak{m}\mathcal{V} = \bigcap_{(v_1, \dots, v_n) \in V_1 \times \dots \times V_n} \text{int}(\text{tconv}(v_1, \dots, v_n)).$$

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